DEMYSTIFYING UNSAFE CODE

Jon Gjengset @jonhoo

OR: "IT'S JUST C" VS "CERTAIN DOOM"

WHEN WE NEED UNSAFE.

WHY UNSAFE?

Lets us write code whose safety relies on **invariants the compiler cannot check**.

Use-cases:

- Working with hardware devices.
- Interacting with external code.
- Writing concurrency primitives.
- Overcoming borrow checker limitations.
- Performance optimizations.

WORKING WITH HARDWARE DEVICES.

Want to show text on screen during boot.

We know things the compiler does not:

- 0xB8000 is mapped to writeable video memory.
- No-one else is writing to 0xB8000.

INTERACTING WITH EXTERNAL CODE.

Want to call into a C library.

Compiler doesn't know if the C code does something unsafe!

We are **asserting** that the compiler can trust the C code.

```
extern "C" { fn c_abs(input: i32) -> i32; }
fn main() {
    println!("{}", unsafe { c_abs(-42)});
}
```

WRITING CONCURRENCY PRIMITIVES.

Want to implement Mutex.

}

Compiler can't check that only **one** &mut T exists **at a time**.

```
fn lock(&self) -> MutexGuard<T> {
   while !self.held.compare_and_swap(false, true, SeqCst) {}
   MutexGuard::new(self)
```

```
impl<T> DerefMut for MutexGuard<T> {
   fn deref_mut(&mut self) -> &mut T {
     unsafe { &mut *self.ptr }
```

OVERCOMING BORROW CHECKER LIMITATIONS.

Want to return reference early.

Compiler will currently reject this valid code.

```
fn next(buffer: &mut String) -> &str {
   loop {
     let event = parse(buffer);
     if true { return event; }
   }}
fn parse(buffer: &mut String) -> &str { ... }
```

PERFORMANCE OPTIMIZATIONS.

Want to remove any and all overheads.

Always measure first. Very rarely worth it.

```
#[repr(C)]
struct SerializedStruct { ... }
```

}

```
unsafe fn cast_deserialize(i: &[u8]) -> &SerializedStruct {
    &*(i.as_ptr() as *const SerializedStruct)
```

UNSAFE GONE WRONG.

WHY NOT UNSAFE?

With unsafe, we tell the compiler that the code is ok, because we checked it manually.

That is, we checked that the code **never** violates **any** part of the Rust type system. We asserting that the unsafe code we wrote **is safe**.

All comes down to: are you sure?

WHY NOT UNSAFE?

The compiler assumes **all** code follows Rust safety rules. That is, it assumes no code **ever**:

- Dereferences dangling/unaligned pointers.
- Violates reference aliasing rules.
- Causes an unsynchronized data race with a write.
- Produces an invalid value.
- ... and a few others.

That includes unsafe code!

EFFECTS OF INCORRECT UNSAFE.

Usually, incorrect unsafe == undefined behavior.

And undefined behavior == russian roulette.

Effect ranges from none to crashes to **arbitrary data corruption**. Effect may not appear today, but appear tomorrow. You have **no guarantees**.

You should **always** avoid undefined behavior.

```
fn compute(i: &u32, o: &mut u32) {
    if *i > 10 {
        *o = 1;
        }
        if *i > 5 {
        *o *= 2;
        }
    }
}
fn compute(i: &u32, o: &mut u32) {
        let cached_i = *i;
        if cached_i > 10 {
            *o = 2;
        } else if cached_i > 5 {
            *o *= 2;
        }
    }
}
```

In Rust, the compiler is allowed to assume that this optimization is okay!

```
impl<T> Vec<T> {
  fn extend_map<U, F>(&mut self, us: &[U], mut f: F)
 where F: FnMut(&U) \rightarrow T {
    self.reserve(us.len()); let cur_len = self.len();
    unsafe { self.set_len(cur_len + us.len()) };
    let into = unsafe { self.as_mut_ptr().add(cur_len) };
    for u in us {
      unsafe { std::ptr::write(into, f(u)) }; into += 1;
    }}}
```

```
What if f() panics?
```

```
impl<T> Drop for LazyDropVec<T> {
  fn drop(&mut self) {
    for v in self.drain() {
        unsafe { COLLECTOR.drop_later(v) };
    }
  }
}
```

What if T contains &mut TcpStream and writes on drop?

```
macro rules! offset of {
  ($t:path, $field:tt) => {
    let uninit = MaybeUninit::<$t>::uninit();
    let ptr = uninit.as_ptr();
    let fptr = unsafe { &(*ptr).$field as *const _ };
    (fptr as usize) - (ptr as usize)
  }
}
```

What if \$t is #[repr(packed)]?

INTEGRITY OF UNSAFE CODE.

Safe interfaces to unsafe code must behave correctly **no matter what the safe code does**.

- Non-deterministic implementation of Eq
- Broken implementation of Ord
- Weird implementations of Deref.

Can only assume **safety** of safe code, not correctness.

INTEGRITY OF "INTERNAL" UNSAFE CODE.

Safe code in the same module can access non-public things!

Encapsulation of unsafe must happen at visibility boundary.

Do not assume callers will remember safety invariants. Never expose unsafe method as safe, even internally!

IS ALL HOPE LOST?

NO!

It is **possible** to write correct unsafe code.

- Be sure you need it.
- Be, like, really sure.
- Read the **nomicon**.
- Be very careful.
- Document **all** unsafe {}
- Run **miri** & ASAN in CI.

AN ASIDE ON MIRI

An interpreter for Rust's mid-level intermediate representation.

Basically, it can run Rust code **in the compiler**.

Can check that the code doesn't do anything "bad".

(works on the playground!)

```
let mut x: Vec<String> = Vec::new();
x.extend_map(&["foo"], |_| panic!());
```

NO!

It is **possible** to write correct unsafe code.

- Be sure you need it.
- Be, like, really sure.
- Read the **nomicon**.
- Be very careful.
- Document **all** unsafe {}
- Run **miri** & ASAN in CI.